

# The categorical logic of elementary arithmetic

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## Computational Logic & Applications

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Au fond de l'Inconnu pour trouver du nouveau!

- On strict extensional reflexivity in compact closed categories  
*Outstanding Contributions in Logic, (to appear March 2023)*  
<https://arxiv.org/abs/2202.08130>
- The inverse semigroup theory of elementary arithmetic  
*(submitted 2022)*  
[www.arxiv.org/abs/2206.07412](http://www.arxiv.org/abs/2206.07412)
- From a conjecture of Collatz to Thompson's group  $F$ , via a conjunction of Girard  
*(submitted 2022)*  
[www.arxiv.org/abs/2202.04443](http://www.arxiv.org/abs/2202.04443)
- Congruential functions via categorical coherence, *(submitted 2022)*
- Operads, Groups, Monoids, & Conjectures. *(Draft manuscript)*

A programming language introduced by John H. Conway (1987)

**Syntax** Programs are (finite?) lists of positive rationals :

$$\left[ \frac{P_0}{Q_0} , \frac{P_1}{Q_1} , \frac{P_2}{Q_2} , \frac{P_3}{Q_3} , \dots \right]$$

**Execution** **Input** is a positive natural number  $n \in \mathbb{N}^+$

**The iterated step :**

- Multiply  $n$  by each  $\frac{P_0}{Q_0} , \frac{P_1}{Q_1} , \frac{P_2}{Q_2} , \frac{P_3}{Q_3} , \dots$  in turn, until a whole number  $n \times \frac{P_j}{Q_j} \in \mathbb{N}^+$  is found.
- replace  $n$  by this number  $n \leftarrow n \times \frac{P_j}{Q_j}$

**Conditional looping** The above step is repeated,  
until the end of the list is reached.

**Output** Either :

- 1 The final value of  $n$ .
- 2 The sequence of values  $n$  takes on during execution.

# An example FRACTRAN program

FRACTRAN is computationally universal

— it can simulate e.g. Turing Machines, or pure untyped  $\lambda$ -calculus

Conway gave the following example :

$$\left[ \begin{array}{cccccccc} \frac{17}{91} & , & \frac{78}{85} & , & \frac{19}{51} & , & \frac{23}{38} & , & \frac{29}{33} & , & \frac{77}{29} & , & \frac{95}{23} & , \\ & & \frac{77}{19} & , & \frac{1}{17} & , & \frac{11}{13} & , & \frac{13}{11} & , & \frac{15}{2} & , & \frac{1}{7} & , & \frac{55}{1} \end{array} \right]$$

On input  $n = 2$ , this program fails to terminate<sup>1</sup>. The powers of two in the infinite sequence generated are :

$$2^2, \dots, 2^3, \dots, 2^5, \dots, 2^7, \dots, 2^{11}, \dots, 2^{13}, \dots, 2^{17}, \dots, 2^{19}, \dots, 2^{23}, \dots, 2^{29}, \dots$$

We recover the list :  $[2^p : p \in \mathbf{Primes}]$ .

**Fun exercise** : What is the action of the program given by i/ *Inverting every fraction*, ii/ *reversing the list*, iii/ *doing both* ??

<sup>1</sup>First proved by Euclid, 350 BCE

# Registers and conditionals from $p$ -adic norms

(The F.T.A.) Every  $n \geq 1$  admits a **unique prime decomposition** :

$$n = 2^{x_0} \times 3^{x_1} \times 5^{x_2} \times 7^{x_3} \times 11^{x_4} \times \dots$$

where a finite number of these  $\{x_j\}_{j \in \mathbb{N}}$  are non-zero. Think of these as **registers**.

Each fraction becomes a *conditional increment* and **branch**

$$\frac{5^3 \times 11^2}{3^3 \times 7^1}$$

COND ( $x_1 \geq 3$  AND  $x_3 \geq 1$ )

$\{x_2 \mapsto x_2 + 3;$   
 $x_4 \mapsto x_4 + 2;$   
BRANCH; }

Conditionals 'consume resources' :

Each (successful) conditional ( $x_j \geq M$ ) decrements  
the **register** by the **test value**  $x_j \mapsto x_j - M$ .

Taking reciprocals interchanges conditionals/decrements with instructions/increments.

# Motivation?

FRACTRAN came out of Conway's work<sup>2</sup> on undecidability in elementary arithmetic

*"Unpredictable Iterations"* – J. H. Conway (1972)

This was demonstrated via the halting problem, based on

**computational universality** of iterative problems on **congruential functions**.

Functions defined "piece-wise linearly on modulo classes"

A function  $f : \mathbb{N} \rightarrow \mathbb{N}$  is **congruential** when there exists a set of modulo classes  $\{A_j\mathbb{N} + B_j\}$  such that

$$f(n) = \frac{X_i n + Y_i}{Z_i} \quad \text{where } n \equiv B_i \pmod{A_i}$$

<sup>2</sup>See also "On E. L. Post's Tag problem" – S. Maslov (1964)

# Some non-trivial motivating examples

- 1 The **Notorious Collatz Conjecture**:  $n \mapsto \begin{cases} \frac{n}{2} & n \text{ even,} \\ \frac{3n+1}{2} & n \text{ odd.} \end{cases}$

(Conjecture: every iterated sequence eventually arrives at 1).

- 2 The **Original Collatz Conjecture** (L. C., unpublished notebooks, 1 July, 1932, as described by J. Lagarias (1985)).

## J. Conway's 'Amusical Permutation'

$$\text{The congruential bijection } \gamma(n) = \begin{cases} \frac{2n}{3} & n \equiv 0 \pmod{3}, \\ \frac{4n-1}{3} & n \equiv 1 \pmod{3}, \\ \frac{4n+1}{3} & n \equiv 2 \pmod{3}, \end{cases}$$

(Conjecture: This function has infinite orbits – the orbit of 8 is  $\infty$ ).

Conway described 2. (the OCC) as,

*"The best candidate we have for a **true but unprovable statement.**"*

# Conjectures as code??

Both the N.C.C. and the O.C.C. are undecided – possibly undecidable<sup>3</sup>.

Can we nevertheless understand them better,  
by interpreting them as `FRACTRAN` programs?

A correspondence of Conway:

Every `FRACTRAN` program implements a **purely multiplicative** congruential function

$$f(n) = \frac{X_i}{Y_i} n + 0 \quad n \equiv B_i \pmod{A_i}$$

... rather neatly ruling out his motivating examples.

**A question:** Instead of looking at simple procedural languages, should we interpret congruential functions via (categorical) logic / lambda calculus??

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<sup>3</sup>Although we could never prove undecidability(!)



# The domain of congruential functions

Congruential functions are defined piece-wise linearly on dissections of  $\mathbb{N}$  into disjoint modulo classes.

## Exact Covering Systems (P. Erdős et al., 1950s - present)

An E.C.S. is an indexed family of modulo classes  $\mathcal{C} = \{A_j\mathbb{N} + B_j\}_{j \in J}$  where:

- $\mathbb{N} = \bigcup_{j \in J} A_j\mathbb{N} + B_j$
- $(A_i\mathbb{N} + B_i) \cap (A_j\mathbb{N} + B_j) = \emptyset \quad \forall i \neq j$

Instead of *indexed sets*, we should take a more *dynamical, algebraic* view, and define them via arrows in a category (monoid).

**A complete triviality** Each modulo class  $A_j\mathbb{N} + B_j$  of an e.c.s. is countably infinite, and so in 1:1 correspondence with  $\mathbb{N}$  itself.

Which arrows exhibit these correspondences,  
& what properties must they satisfy?

# Dynamical Algebras via exact covering systems (I)

Let  $\mathcal{C} = \{A_i\mathbb{N} + B_i\}_{i=1}^n$  be an exact covering system

- 1  $A_i\mathbb{N} + B_i$  is bijective with  $\mathbb{N}$  for all  $i = 1..n$

This is exhibited by the injection :  $p_i(n) \stackrel{\text{def.}}{=} A_i n + B_i$

- 2 Elements of  $\mathcal{C}$  are pairwise disjoint

becomes the condition :  $im(p_i) \cap im(p_j) = \emptyset \quad \forall i \neq j$

- 3 Elements of  $\mathcal{C}$  cover the whole of  $\mathbb{N}$

is equivalent to :  $\bigcup_{i=1}^n im(p_i) = \mathbb{N}$

# Dynamical Algebras via exact covering systems (II)

Instead of discussing *domains* and *images* ...

We work with the 'symmetric inverse monoid'  $\mathcal{I}(\mathbb{N})$  of **partial injective functions**

$$f(x) = f(x') \Rightarrow x = x' \text{ provided both are defined.}$$

and their unique **generalised inverses**  $f(x) = y \Leftrightarrow f^*(y) = x$

The generalised inverse of  $p_j$  is given by

$$p_j^*(n) = \begin{cases} \frac{n-B_j}{A_j} & n \equiv B_j \pmod{A_j} \\ \perp & \text{otherwise} \end{cases}$$

The required conditions are the defining relations of the  $n$ -th **dynamical algebra** :

$$p_j^* p_i = \begin{cases} Id & i = j \\ \emptyset & i \neq j \end{cases} \quad \text{and} \quad \sum_{i=1}^n p_i p_i^* = Id$$

where 'sum' is union of partial functions with disjoint domains / images.

(Categorically, "sum" is the join w.r.t. the poset-enrichment of the category of partial injective functions).

# Congruential functions in logic & lambda calculi

Given an exact covering system  $\{A_i\mathbb{N} + B_i\}_{i=0}^{n-1}$ , the corresponding partial injections

$$p_0, p_1, \dots, p_{n-1} : \mathbb{N} \rightarrow \mathbb{N}$$

generate a copy of the  $n$ -th **dynamical algebra** (a.k.a. Nivat & Perot's  $n$ -th polycyclic monoid), as (partial) congruential functions.

These appear in a series of (closely related) logical / computational models :

**M.E.L.L. and system  $\mathcal{F}$**  *Geometry of Interaction I and II*

– J.-Y. Girard (1988,1990),

**untyped lambda calculus** *Local and asynchronous  $\beta$ -reduction*

– V. Danos & L. Regnier (1992),

**combinatory logic** *Gol and linear combinatory algebra*

– S. Abramsky, E. Hagverdi, P. Scott (2002),

...

The examples that motivated Conway arise when  
we consider the **category theory** behind these models.

# A categorically significant example

In J.-Y. Girard's Geometry of Interaction system (Parts I, II) :

**Propositions** are modelled by partial injective functions – elements of  $\mathcal{I}(\mathbb{N})$ .

**Conjunction** is defined in terms of the two-generator dynamical algebra, based on the odd-even dissection  $\mathbb{N} = 2\mathbb{N} \cup 2\mathbb{N} + 1$ .

Expanding out his definition, we get :

$$(f \star g)(2n) = 2.f(n)$$

$$(f \star g)(2n + 1) = 2.g(n) + 1$$

A simple description, based on Hilbert's Grand Hotel

This “writes two functions as a single function”, by

*replicating their behaviour on the **even** and **odd** numbers respectively.*

**Key point** : Girard's conjunction  $(- \star -) : \mathcal{I}(\mathbb{N}) \times \mathcal{I}(\mathbb{N}) \leftrightarrow \mathcal{I}(\mathbb{N})$  is a (semi-monoidal) categorical tensor on a monoid.

# Congruential Canonical Isomorphisms

Girard's conjunction cannot be strictly associative

Strict associativity of a faithful (injective) tensor on a monoid



The unique object of the monoid is the unit object.

*Coherence & Strictification for Self-Similarity*

Journal Homotopy & Related Structures (P.M.H. 2016)

There is a non-trivial natural isomorphism  $(- \star (- \star -)) \Rightarrow ((- \star -) \star -)$

$$\alpha(a \star (b \star c)) = ((a \star b) \star c)\alpha \quad \forall a, b, c \in \mathcal{I}(\mathbb{N})$$

whose unique component (the **associator**)  $\alpha(n) = \begin{cases} 2n & n \equiv 0 \pmod{2}, \\ n + 1 & n \equiv 1 \pmod{4}, \\ \frac{n-1}{2} & n \equiv 3 \pmod{4}, \end{cases}$

is a congruential function, satisfying MacLane's pentagon condition

$$\alpha^2 = (\alpha \star Id) \alpha (Id \star \alpha)$$

# Collatz-like problems for canonical isomorphisms?

## Question :

Can we decide whether orbits of natural numbers under the associativity isomorphism  $\alpha$  are **finite** or **infinite**?

The associator  $\alpha$  has disappointingly simple behaviour under iteration :

$n$  is even  $\alpha(n) = 2n > n$  for all  $n \neq 0$

$n$  is odd Two possibilities :

- 1  $\alpha(n)$  is odd  $\Rightarrow \alpha(n) < n$
- 2  $\alpha(n)$  is even.

All orbits, apart from the fixed point  $\alpha(0) = 0$ , are infinite.

We see similar *disturbingly trivial* behaviour under iteration for arbitrary associativity isomorphisms<sup>4</sup>

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<sup>4</sup>i.e. components of natural isomorphisms between  $k$ -fold bracketings of Girard's conjunction.

# Briefing for a descent into group theory

The structure group for the associativity identity, P. Dehornoy (1996)

Given a tensor  $\alpha$  on a monoid, the associativity isomorphisms form a group, isomorphic to R. Thompson's group  $\mathcal{F} = \langle X_j : X_j X_k X_j^{-1} = X_{k-1} \quad \forall j < k - 1 \in \mathbb{N} \rangle$ .

A suitable generating set is given by

$$X_0 = \alpha, X_1 = (Id \star \alpha), X_2 = Id \star (Id \star \alpha), \dots$$

since

$$\alpha(Id \star (Id \star X_j))\alpha^{-1} = (Id \star Id) \star \alpha = Id \star \alpha$$

$X_0, X_1, X_2, \dots$  can be thought of as: “ $\alpha$  replicated on

$$2\mathbb{N} + 1, 4\mathbb{N} + 3, 8\mathbb{N} + 7, \dots$$

and the identity elsewhere.”

Again, all orbits are either **infinite** or **fixed points**.



# Characterising canonical isomorphisms for Girard's conjunction

## A standard result

The above 'infinite presentation' of Thompson's  $\mathcal{F}$  is not minimal;  
 $X_0$  and  $X_1$  suffice to generate the whole group.

**Corollary** The two bijections

$$\alpha(n) = \begin{cases} 2n & n \pmod{2} = 0 \\ n+1 & n \pmod{4} = 1 \\ \frac{n-1}{2} & n \pmod{4} = 3 \end{cases} \quad (Id \star \alpha)(n) = \begin{cases} n & n \pmod{2} = 0 \\ 2n-1 & n \pmod{4} = 1 \\ n+2 & n \pmod{8} = 3 \\ \frac{n-1}{2} & n \pmod{8} = 7 \end{cases}$$

generate a copy of Thompson's group  $\mathcal{F}$ , and so capture *all* canonical associativity isomorphisms for  $\_ \star \_$ .

None of these have interesting behaviour under iteration!

# An important question(!)

Are congruential canonical isomorphisms always so uninteresting<sup>5</sup>??

“Given any  $n$  objects of a monoidal category, the canonical associativity isomorphisms give a commuting diagram whose shape is the 1-skeleton of the  $n^{\text{th}}$  associahedron  $\mathcal{K}_n$ ”. — M. Kapranov (1993)

## The claim :

We see **non-trivial behaviour**<sup>a</sup> when we move off the 1-skeleton, and consider commuting diagrams of maps between more general facets of associahedra.

<sup>a</sup>Precisely, the operators of Collatz that motivated Conway . . .

<sup>5</sup>from an *iterative*, rather than *group-theoretic* viewpoint!

# An unbiased series of conjunctions

Girard gave a **binary** model of conjunction  $(\_ \star \_ ) : \mathcal{I}(\mathbb{N}) \times \mathcal{I}(\mathbb{N}) \leftrightarrow \mathcal{I}(\mathbb{N})$ .

“( $a \star b$ ) replicates  $a, b$  on the congruence classes  $2\mathbb{N}, 2\mathbb{N} + 1$  respectively”.

- We draw this as



There is an obvious **ternary** analogue,  $(\_ \star \_ \star \_ ) : \mathcal{I}(\mathbb{N}) \times \mathcal{I}(\mathbb{N}) \times \mathcal{I}(\mathbb{N}) \leftrightarrow \mathcal{I}(\mathbb{N})$

$$(a \star b \star c)(3n + i) = \begin{cases} 3.a(n) & i = 0 \\ 3.b(n) + 1 & i = 1 \\ 3.c(n) + 2 & i = 2 \end{cases}$$

“Replicate  $a, b, c$  on the congruence classes  $3\mathbb{N}$  ,  $3\mathbb{N} + 1$  ,  $3\mathbb{N} + 2$  respectively”.

- We draw this as



# The general case :

For any  $k \geq 1$ , we form the  $k^{\text{th}}$  **unbiased conjunction** by :

$$(f_0 \star \dots \star f_{k-1})(kn + i) = k \cdot f_i(n) + i \text{ where } i = 0, 1, 2, \dots, k - 1$$

Alternatively & equivalently,

$$(f_0 \star \dots \star f_{k-1})(x) = k \cdot f_i\left(\frac{x-i}{k}\right) + i \text{ where } x \equiv i \pmod k$$

This gives, for any  $k > 0$ , an injective homomorphism  $\mathcal{I}(\mathbb{N})^{\times k} \hookrightarrow \mathcal{I}(\mathbb{N})$  that :

*replicates  $f_0, f_1, \dots, f_{k-1}$  on the congruence classes modulo  $k$*

*i.e.  $\{k\mathbb{N}, k\mathbb{N} + 1, \dots, k\mathbb{N} + (k - 1)\}$*

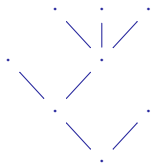
For  $k = 1, 2, 3, 4, \dots$ , we draw the unbiased conjunctions as



We may compose these, to build an **operad**.

# Composing elementary conjunctions

These 'compose by substitution' to give an operad  $\mathcal{GC}$  of **generalised conjunctions**. Each  $k$ -leaf tree determines an injective hom.  $\mathcal{I}(\mathbb{N})^{\times k} \hookrightarrow \mathcal{I}(\mathbb{N})$ .



$$: \mathcal{I}(\mathbb{N})^{\times 5} \hookrightarrow \mathcal{I}(\mathbb{N})$$

$$(f_0, f_1, f_2, f_3, f_4) \mapsto ((f_0 \star (f_1 \star f_2 \star f_3)) \star f_4)$$

More formally :

Consider the (non-symmetric) endomorphism operad of  $\mathcal{I}(\mathbb{N})$  within the category  $(\mathbf{Inv}, \times)$  of inverse monoids, with Cartesian product. This contains one *unbiased conjunction* of each arity  $> 0$ .

These generate the sub-operad  $\mathcal{GC}$  of generalised conjunctions.

# A freely generated operad

**Claim :** The operad  $\mathcal{GC}$  is isomorphic to the operad  $\mathcal{RPT}$  of **rooted planar trees**.

Thus

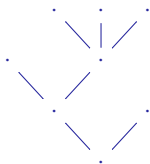
- It is *freely* generated by the unbiased conjunctions.
- Every distinct rooted planar tree determines a distinct generalised conjunction.
- The  $n$ -ary operations  $\mathcal{GC}_n$  are in 1:1 correspondence with facets of the  $n^{\text{th}}$  associahedron  $\mathcal{K}_n$ .

To be proved :

Each tree in  $\mathcal{RPT}_k$  determines a *distinct* homomorphism  $\mathcal{I}(\mathbb{N})^{\times k} \hookrightarrow \mathcal{I}(\mathbb{N})$ .

# Back to simple arithmetic!

The tree



defines a homomorphism :  $\mathcal{I}(\mathbb{N})^{\times 5} \hookrightarrow \mathcal{I}(\mathbb{N})$

In the generalised conjunction  $((f_0 \star (f_1 \star f_2 \star f_3)) \star f_4)$ , the action of each  $f_j$  is mapped :

from The whole of the natural numbers  $\mathbb{N}$

to Some modulo class  $A_j\mathbb{N} + B_j$ .

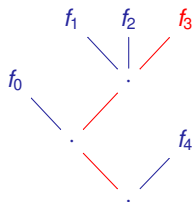
**For example :**  $f_3$  is 'replicated' on the modulo class  $12\mathbb{N} + 10$ .

Obvious Question:

How do we derive these coefficients *from the tree*?

# A root and branch approach

Deriving  $12\mathbb{N} + 10$ , from the leaf-to-root path :



Branch number 2 of 3

Branch number 1 of 2

Branch number 0 of 2

Multiplicative coefficient :  $12 = 3 \times 2 \times 2$

Additive coefficient : 

(Decimal)	=	<u>Base 3</u>	=	<u>Base 2</u>	=	<u>Base 2</u>
10		2		1		0

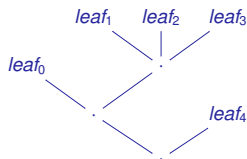
Positional **mixed-radix** number systems

First formal study by G. Cantor, *Über einfache Zahlensysteme* (1869)



# Covering the Numbers with Trees

Rooted Planar Trees are uniquely determined by the addresses of their leaves



$leaf_0$		$(0, 2)$	$(0, 2)$	$4\mathbb{N}$
$leaf_1$	$(0, 3)$	$(1, 2)$	$(0, 2)$	$12\mathbb{N} + 2$
$leaf_2$	$(1, 3)$	$(1, 2)$	$(0, 2)$	$12\mathbb{N} + 6$
$leaf_3$	$(2, 3)$	$(1, 2)$	$(0, 2)$	$12\mathbb{N} + 10$
$leaf_4$		$(1, 2)$		$2\mathbb{N} + 1$

which uniquely determine ordered exact covering systems, such as

$$4\mathbb{N} , 12\mathbb{N} + 2 , 12\mathbb{N} + 6 , 12\mathbb{N} + 10 , 2\mathbb{N} + 1$$

## Corollaries:

- 1 Distinct trees determine distinct homomorphisms, so  $\mathcal{GC} \cong \text{RPT}$ .
- 2 Distinct  $k$ -leaf trees determine distinct realisations of the  $k$ -th dynamical algebra.

# Mapping between generalised conjunctions

The operations of  $\mathcal{GC}_k$  are homomorphisms (functors).

Given generalised conjunctions  $T, U : \mathcal{I}(\mathbb{N})^{\times k} \hookrightarrow \mathcal{I}(\mathbb{N})$ , how can we find (well-behaved) natural isomorphisms between them?

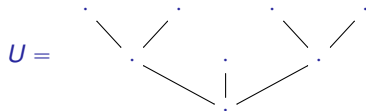
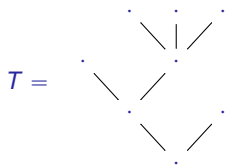
$$\begin{array}{ccc} & T & \\ & \curvearrowright & \\ \mathcal{I}(\mathbb{N})^{\times k} & & \mathcal{I}(\mathbb{N}) \\ & \Downarrow ?? & \\ & \curvearrowleft & \\ & U & \end{array}$$

**Convention :** As generalised conjunctions are **monoid** homomorphisms, natural transformations have a **single** component.

We identify nat. iso.s with their unique component in  $\mathcal{I}(\mathbb{N})$ .

# Congruential functions as natural isomorphisms

Consider the generalised conjunctions<sup>6</sup>  $T, U : \mathcal{I}(\mathbb{N})^{\times 5} \leftrightarrow \mathcal{I}(\mathbb{N})$



We build a natural isomorphism  $\eta_{T,U} : T \Rightarrow U$  by monotonically mapping between their respective ordered exact covering systems :

<b>leaf 0</b>	$4\mathbb{N}$	$\mapsto$	$6\mathbb{N}$
<b>leaf 1</b>	$12\mathbb{N} + 2$	$\mapsto$	$6\mathbb{N} + 3$
<b>leaf 2</b>	$12\mathbb{N} + 6$	$\mapsto$	$3\mathbb{N} + 1$
<b>leaf 3</b>	$12\mathbb{N} + 10$	$\mapsto$	$6\mathbb{N} + 2$
<b>leaf 4</b>	$2\mathbb{N} + 1$	$\mapsto$	$6\mathbb{N} + 4$

This gives, as desired,

$$\eta_{T,U}.((a \star (b \star c \star d)) \star e) = ((a \star b) \star c \star (d \star e)).\eta_{T,U}$$

<sup>6</sup>edges of the fifth associahedron  $\mathcal{K}_5$

# Concrete formulæ for congruential natural isomorphisms

Given two ordered exact covering systems, determined by  $k$ -ary generalised conjunctions  $T, U$

$$\text{leaf } 0 \quad A_0\mathbb{N} + B_0 \quad \mapsto \quad C_0\mathbb{N} + D_0$$

$$\text{leaf } 1 \quad A_1\mathbb{N} + B_1 \quad \mapsto \quad C_1\mathbb{N} + D_1$$

$$\vdots \quad \quad \quad \vdots \quad \quad \quad \vdots$$

$$\text{leaf } k - 1 \quad A_{k-1}\mathbb{N} + B_{k-1} \quad \mapsto \quad C_{k-1}\mathbb{N} + D_{k-1}$$

The natural isomorphism  $\eta_{T,U}$  is the congruential bijection

$$\eta_{T,U}(x) = \frac{1}{A_j} \left( C_j x + \begin{vmatrix} A_j & B_j \\ C_j & D_j \end{vmatrix} \right) \quad \text{where } x \equiv B_j \pmod{A_j}$$

Think of these as, “mapping between distinct realisations of the  $k$ -th dynamical algebra”.

## Some context :

These are a *special kind* of congruential function, where

*each element of an e.c.s. is mapped linearly to another modulo class.*

- **A computational interpretation :**

Functional equations associated with congruential functions,  
Theoretical Computer Science *S. Berckel (1994)*

- **Pure group theory :**

A simple group given by interchanging residue classes of the integers  
Mathematische Zeitschrift *S. Kohl (2010)*

- **Close connections with the N.C.C. :**

The  $[3x + 1]$  Collatz conjecture in a group-theoretic context  
Journal of Group Theory *S. Kohl (2017)*

- **Normal forms & formulæ for composition :**

The inverse semigroup theory of elementary arithmetic  
[www.arxiv.org/2206.07412](http://www.arxiv.org/2206.07412) *P.M.H. (2022)*

# A groupoid of generalised conjunctions

Observe that :

- $\eta_{T,T} = Id \in \mathcal{I}(\mathbb{N})$
- $\eta_{T,U} \eta_{S,T} = \eta_{S,U}$
- $\eta_{T,U}^{-1} = \eta_{U,T}$

We have a **posetal groupoid**<sup>7</sup>  $\mathcal{C}\omega$  of functors / natural iso.s, given by :

**Objects** Generalised conjunctions (operations of  $\mathcal{GC}$ )

**Arrows**  $\mathcal{C}\omega(T, U) = \begin{cases} \{\eta_{T,U}\} & T, U \text{ have the same arity,} \\ \emptyset & \text{otherwise.} \end{cases}$

A reminder :

**Arrows** are natural isomorphisms between homomorphisms  $\mathcal{I}(\mathbb{N})^{\times k} \hookrightarrow \mathcal{I}(\mathbb{N})$  determined by their unique components in  $\mathcal{I}(\mathbb{N})$

<sup>7</sup>within which, 'all diagrams commute'

# Generalised conjunctions within a posetal category

**Claim:** As well as being the **objects** of the posetal groupoid  $\mathcal{C}\omega$ , generalised conjunctions are also **functors**  $\prod^k \mathcal{C}\omega \rightarrow \mathcal{C}\omega$ .

**On Objects :** this is defined by operadic composition

$$(T_0 \star \dots \star T_x) \stackrel{\text{def.}}{=} \begin{array}{c} T_0 \quad T_1 \quad \dots \quad T_x \\ \searrow \quad \searrow \quad \quad \quad \searrow \\ \cdot \end{array}$$

**On Arrows :** We take the generalised conjunction of unique components.

To check **functoriality** :

$$\eta_{(T_0 \star \dots \star T_x), (U_0 \star \dots \star U_x)} = (\eta_{T_0, U_0} \star \dots \star \eta_{T_x, U_x})$$

$$\begin{array}{ccccccc} T_0 & & T_1 & & \dots & & T_k \\ \parallel & & \parallel & & & & \parallel \\ \eta_{T_0, U_0} \downarrow & & \eta_{T_1, U_1} \downarrow & & & & \eta_{T_k, U_k} \downarrow \\ U_0 & & U_1 & & \dots & & U_x \end{array}$$

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$$\begin{array}{c} T_0 \quad T_1 \quad \dots \quad T_x \\ \diagdown \quad \diagup \quad \dots \quad \diagdown \quad \diagup \\ \cdot \end{array} \xrightarrow{\eta_{T_0, U_0} \star \dots \star \eta_{T_x, U_x}} \begin{array}{c} U_0 \quad U_1 \quad \dots \quad U_x \\ \diagdown \quad \diagup \quad \dots \quad \diagdown \quad \diagup \\ \cdot \end{array}$$



# Operads all the way down

## Key points :

- The groupoid  $\mathcal{C}\omega$  is **posetal**, with objects in 1:1 correspondence with facets of associahedra.
- There is a **unique arrow** between any two facets of the same associahedron.
- It admits an  $\mathbb{N}^+$ -indexed family of faithful functors :

$$\begin{aligned} Id & : \mathcal{C}\omega \rightarrow \mathcal{C}\omega \\ (- \star -) & : \mathcal{C}\omega \times \mathcal{C}\omega \rightarrow \mathcal{C}\omega \\ (- \star - \star -) & : \mathcal{C}\omega \times \mathcal{C}\omega \times \mathcal{C}\omega \rightarrow \mathcal{C}\omega \\ & \vdots \end{aligned}$$

- As a **very special case**, it contains a copy of MacLane's posetal monoidal groupoid  $(\mathcal{W}, \square)$ .

## Fun question :

What does the endomorphism operad<sup>a</sup> of  $\mathcal{C}\omega$  look like, in the category  $(\mathbf{Cat}, \times)$  ?

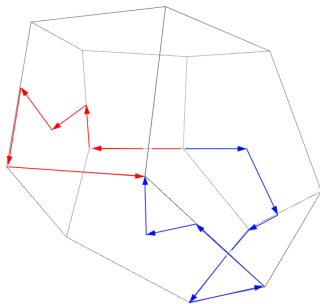
<sup>a</sup>or functor category thereof ...

# Commuting diagrams from associahedra

As  $\mathcal{C}_\omega$  is posetal, all diagrams over it commute.

There exists a unique arrow between any two objects  
iff  
they are facets of the same associahedron.

We may build commuting diagrams from paths between arbitrary facets.



Labels on edges are **congruential functions**.

# Concrete examples : the third associahedron

$$\begin{array}{ccc}
 ((- \star -) \star -) & \xleftarrow{\alpha} & (- \star (- \star -)) \\
 & \searrow^{\gamma_b} & \swarrow_{\gamma} \\
 & (- \star - \star -) &
 \end{array}$$

$$\alpha(n) = \begin{cases} 2n & n \equiv 0 \pmod{2}, \\ n+1 & n \equiv 1 \pmod{4}, \\ \frac{n-1}{2} & n \equiv 3 \pmod{4}. \end{cases}$$

$$\gamma(n) = \begin{cases} \frac{2n}{3} & n \equiv 0 \pmod{3}, \\ \frac{4n-1}{3} & n \equiv 1 \pmod{3}, \\ \frac{4n+1}{3} & n \equiv 2 \pmod{3}. \end{cases}$$

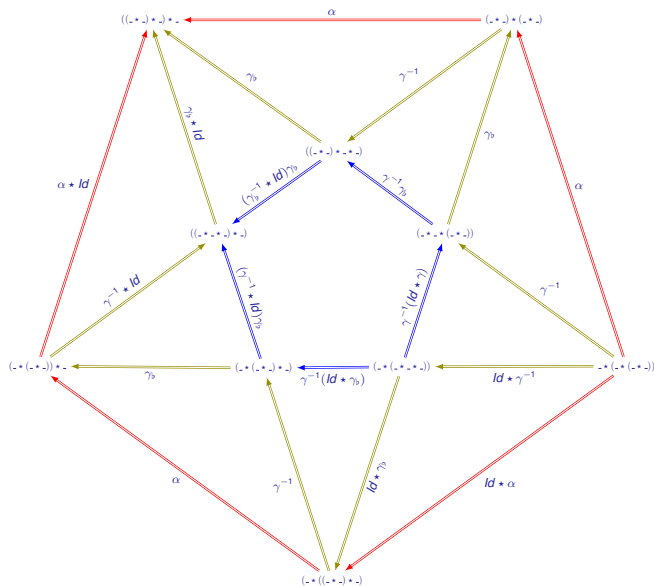
$$\gamma_b(n) = \begin{cases} \frac{4n}{3} & n \equiv 0 \pmod{3}, \\ \frac{4n+2}{3} & n \equiv 1 \pmod{3}, \\ \frac{2n-1}{3} & n \equiv 2 \pmod{3}. \end{cases}$$

## We have already seen these (!)

- $\alpha$  : the **associativity isomorphism** for Girard's conjunction.
- $\gamma$  : the '**amusical permutation**', from the Original Collatz Conjecture.
- $\gamma_b$  related to  $\gamma$  by a natural transformation :  $1 + \gamma_b(n) = \gamma(n+1)$ .

This '*replicates the behaviour of  $\gamma$ , one step down*'.

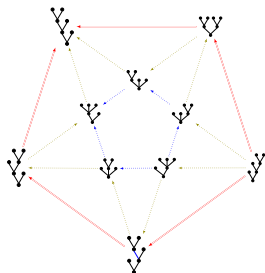
# From the third to the fourth associahedron



# A Convergent Series?

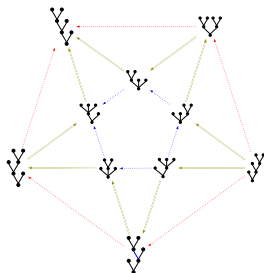
MacLane's pentagon is the 1-skeleton of Stasheff's associahedron  $\mathcal{K}_4$ ;  
we understand the rest of the pentagram in similar terms.

Mapping between Vertices



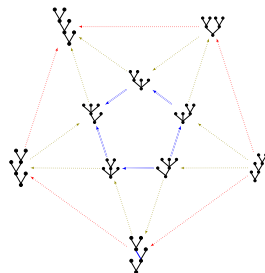
The associator for Girard's  $(- \star -)$

Mapping between Edges & Vertices



Collatz's bijection(s)

Mapping between Edges



Girard-Collatz composites

These are natural isomorphisms in a posetal functor category  
and also, congruential bijections on the natural numbers.

# From associativity to symmetry

One of Conway's motivating examples (the O.C.C.) arises from :

- Coherence for **associativity**
- Re-bracketing generally

The other motivating example (the N.C.C.) is seen when we extend to :

- Coherence for **symmetry & associativity**,
- Non-symmetric operads.

The core operator from the  $3x + 1$  problem is not a bijection!

The problem has, however, been characterised in terms of (congruential) bijections

**Claim** : these are coherence isomorphisms

## Journal of Group Theory, Stefan Kohl (2017)

“The [Notorious] Collatz conjecture in a group-theoretic context”

Many thanks to Matt Brin (Binghamton) for bringing this to my attention!

Introduces a three-generator group of permutations  $\mathcal{G}_{\mathcal{T}} \subseteq \mathcal{I}(\mathbb{N})$  satisfying :

“ $\mathcal{G}_{\mathcal{T}}$  acts transitively<sup>8</sup> on  $\mathbb{N}$  iff the  $3x + 1$  conjecture is true”.

The generators – and hence all elements of  $\mathcal{G}_{\mathcal{T}}$  – are **congruential bijections**.

**Claim :** Once we include symmetry, these are natural isomorphisms between generalised conjunctions, describing coherence for  $(- \star -)$  and  $(- \star - \star -)$ .

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<sup>8</sup>For all  $x, y \in \mathbb{N}$ , there exists  $f \in \mathcal{G}_{\mathcal{T}}$  such that  $f(x) = y$ .

# What has been missing . . .

Girard's conjunction  $_ * _ : \mathcal{I}(\mathbb{N}) \times \mathcal{I}(\mathbb{N}) \rightarrow \mathcal{I}(\mathbb{N})$  is a **symmetric** tensor.

$$\sigma(f * g) = (g * f)\sigma \quad \forall g, f \in \mathcal{I}(\mathbb{N})$$

where the **symmetry iso.**  $\sigma \in \mathcal{I}(\mathbb{N})$  is given by :  $\sigma(n) = \begin{cases} n + 1 & n \text{ even,} \\ n - 1 & n \text{ odd.} \end{cases}$

*Coherence for symmetry* requires **MacLane's hexagon condition** :

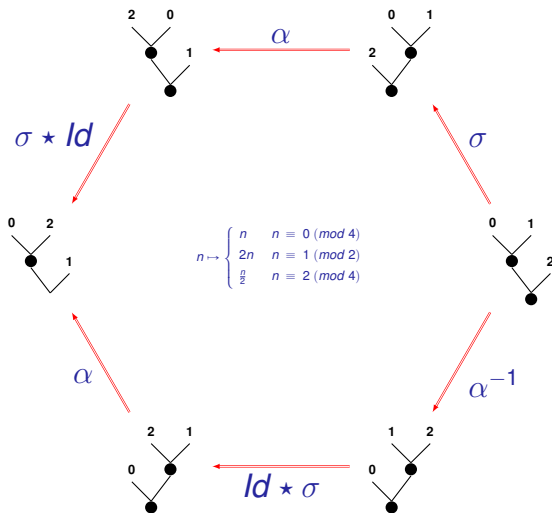
$$\alpha\sigma\alpha = (\sigma * Id)\alpha(Id * \sigma)$$

which arises from the 1-skeleton of Kapranov's third **permutoassociahedron**,  $\mathcal{KP}_3$ , based on *symmetric operads* with an action of  $\mathcal{S}(\{0, \dots, k - 1\})$  on leaves of trees.



# MacLane's hexagon for Girard's conjunction

Vertices are **bracketings** of **orderings** of three symbols,  $\{0, 1, 2\}$ .



# The generators of $\mathcal{G}_{\mathcal{T}}$

From  $\mathcal{KP}_2$  The symmetry isomorphism for Girard's conjunction,  $\sigma \in \mathcal{I}(\mathbb{N})$

$$\begin{array}{c} 1 \quad 0 \\ \diagdown \quad / \\ \bullet \end{array} \quad \Leftarrow \quad \begin{array}{c} 0 \quad 1 \\ \diagdown \quad / \\ \bullet \end{array} \quad n \mapsto \begin{cases} n+1 & n \equiv 0 \pmod{2} \\ n-1 & n \equiv 1 \pmod{2} \end{cases}$$

From  $\mathcal{KP}_3$  The bijection from MacLane's hexagon  $(\sigma \star Id)\alpha\sigma = \alpha(Id \star \sigma)\alpha^{-1}$

$$\begin{array}{c} 0 \quad 2 \\ \diagdown \quad / \\ \bullet \quad \diagdown \quad / \\ \bullet \quad 1 \end{array} \quad \Leftarrow \quad \begin{array}{c} 0 \quad 1 \\ \diagdown \quad / \\ \bullet \quad \diagdown \quad / \\ \bullet \quad 2 \end{array} \quad n \mapsto \begin{cases} n & n \equiv 0 \pmod{4} \\ 2n & n \equiv 1 \pmod{2} \\ \frac{n}{2} & n \equiv 2 \pmod{4} \end{cases}$$

From  $\mathcal{KP}_5$  A mapping between *edges* of the fifth<sup>9</sup> permutaoassiahedron.

$$\begin{array}{c} 0 \quad 3 \quad 2 \quad 1 \quad 4 \\ \diagdown \quad / \quad \diagdown \quad / \\ \bullet \quad \bullet \quad \diagdown \quad / \\ \bullet \end{array} \quad \Leftarrow \quad \begin{array}{c} 0 \quad 1 \quad 2 \quad 3 \quad 4 \\ \diagdown \quad / \quad \diagdown \quad / \\ \bullet \quad \bullet \quad \diagdown \quad / \\ \bullet \end{array} \quad n \mapsto \begin{cases} n & n \equiv 0 \pmod{6} \\ \frac{2n+2}{3} & n \equiv 2 \pmod{6} \\ n & n \equiv 4 \pmod{6} \\ \frac{3n-1}{2} & n \equiv 1 \pmod{4} \\ n & n \equiv 2 \pmod{4} \end{cases}$$

*All of these decompose into composites & generalised conjunctions of  $\alpha, \sigma, \gamma, \gamma_b$ .*

<sup>9</sup>Caution :  $\mathcal{KP}_n$  is not **simple**, for  $n \geq 4$ .

# That third generator, via canonical isomorphisms

Using **vertex-to-vertex** and **vertex-to-edge** maps, we implement the required **edge-to-edge** map.

