

Counting Environments and Closures

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1 What is a closure ?

2 How to count closures?

3 Plain

4 Closed

5 Conclusion

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 - ★ If one “reduce” this closure, one gets $\lambda \lambda \underline{0}$.

The equations for closures

The definition of the set of closures and of the set of environments is recursive

$$\mathit{Clos} ::= \langle \Lambda, \mathit{Env} \rangle$$

$$\mathit{Env} ::= \square \mid \mathit{Clos} : \mathit{Env}$$

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- λ -terms:

$$|0| = 1$$

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$$|M N| = |M| + |N| + 1$$

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- Closures and environments:

$$|\langle M, e \rangle| = |M| + |e| \qquad |c : e| = |c| + |e| \qquad |\square| = 0.$$

The first environments and closures

Here are the first **plain environments**:

size	environments	total
0	\square	1
1	$\langle \underline{0}, \square \rangle : \square$	1
2	$\langle \underline{0}, \square \rangle : \langle \underline{0}, \square \rangle : \square$ $\langle \underline{0}, \langle \underline{0}, \square \rangle : \square \rangle : \square$ $\langle \lambda \underline{0}, \square \rangle : \square, \quad \langle \underline{1}, \square \rangle : \square$	4

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and **plain closures**:

size	closures	total
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1	$\langle \underline{0}, \square \rangle$	1
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▶ Closed Closures

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Equations for plain closures and environments

- Set equations

$$\begin{aligned}\mathcal{E}_\infty &= \mathcal{C}_\infty : \mathcal{E}_\infty \mid \square \\ \mathcal{C}_\infty &= \langle \mathcal{L}_\infty, \mathcal{E}_\infty \rangle.\end{aligned}$$

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- Equations for generating functions

$$\begin{aligned}E_\infty(z) &= C_\infty(z)E_\infty(z) + 1 \\ C_\infty(z) &= L_\infty(z)E_\infty(z).\end{aligned}$$

Generating functions

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$$E_{\infty}(z) = \frac{1 - \sqrt{1 - 4L_{\infty}(z)}}{2L_{\infty}(z)}$$
$$C_{\infty}(z) = \frac{1}{2} \left(1 - \sqrt{1 - 4L_{\infty}(z)} \right).$$

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$$\rho < \rho_{\infty}$$

- Radius of convergence in 0

$$\rho = \frac{1}{10} \left(25 - \sqrt{545} \right) \doteq 0.165476.$$

- Newton-Puiseux expansions

$$E_{\infty}(z) = 2 + -\frac{1}{4} \sqrt{\frac{5}{47} (109 + 35\sqrt{545})} \sqrt{1 - \frac{z}{\rho}} + O\left(\left|1 - \frac{z}{\rho}\right|\right)$$

$$C_{\infty}(z) = \frac{1}{2} + \frac{2 \sqrt{\frac{10(48069\sqrt{5} - 10295\sqrt{109})}{65\sqrt{109} - 301\sqrt{5}}}}{3\sqrt{545} - 77} \sqrt{1 - \frac{z}{\rho}} + O\left(\left|1 - \frac{z}{\rho}\right|\right)$$

Asymptotics

The numbers e_n and c_n of plain environments and closures of size n , respectively, admit the following asymptotic approximations:

$$e_n \sim C_e \cdot \rho^{-n} n^{-3/2}$$

$$c_n \sim C_c \cdot \rho^{-n} n^{-3/2}$$

where

$$C_e = \frac{\sqrt{\frac{5}{47} (109 + 35\sqrt{545})}}{8\sqrt{\pi}} \doteq 0.699997,$$

$$C_c = \frac{\sqrt{\frac{10(48069\sqrt{5} - 10295\sqrt{109})}{65\sqrt{109} - 301\sqrt{5}}}}{\sqrt{\pi} (77 - 3\sqrt{545})} \doteq 0.174999$$

and

$$\rho = \frac{1}{10} (25 - \sqrt{545}) \doteq 0.165476 \quad \text{giving} \quad \rho^{-n} \doteq 6.04315^n.$$

- The generating functions are **algebraic**

Holonomicity

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- then we get **holonomic functions**
- then, using **gfun** Maple package, we get **linear recurrences**:

Linear recurrences

$$\begin{aligned} & (125 n^3 - 125 n) e_n + \\ & (-475 n^3 - 150 n^2 + 325 n) e_{n+1} + \\ & (-1625 n^3 - 13650 n^2 - 29125 n - 17100) e_{n+2} + \\ & (5925 n^3 + 65550 n^2 + 204825 n + 190800) e_{n+3} + \\ & (-10950 n^3 - 149850 n^2 - 609000 n - 744300) e_{n+4} + \\ & (43599 n^3 + 638460 n^2 + 3028701 n + 4633680) e_{n+5} + \\ & (-97781 n^3 - 1680378 n^2 - 9481237 n - 17550960) e_{n+6} + \\ & (122749 n^3 + 2388066 n^2 + 15211685 n + 31648968) e_{n+7} + \\ & (-184402 n^3 - 3954630 n^2 - 27717140 n - 63149544) e_{n+8} + \\ & (280081 n^3 + 6826380 n^2 + 54868451 n + 145130568) e_{n+9} + \\ & (-205649 n^3 - 5654610 n^2 - 51851989 n - 158722620) e_{n+10} + \\ & (37439 n^3 + 1339686 n^2 + 16635271 n + 70682784) e_{n+11} + \\ & (-68686 n^3 - 3028038 n^2 - 43616336 n - 205972920) e_{n+12} + \\ & (222029 n^3 + 9258780 n^2 + 128417911 n + 592399800) e_{n+13} + \\ & (-241115 n^3 - 10519830 n^2 - 152823475 n - 739190880) e_{n+14} + \\ & (134151 n^3 + 6201222 n^2 + 95476551 n + 489605640) e_{n+15} + \\ & (-42231 n^3 - 2067834 n^2 - 33729375 n - 183277332) e_{n+16} + \\ & (7470 n^3 + 386418 n^2 + 6659316 n + 38233296) e_{n+17} + \\ & (-678 n^3 - 36972 n^2 - 671670 n - 4065240) e_{n+18} + \\ & (24 n^3 + 1380 n^2 + 26436 n + 168720) e_{n+19} = 0. \end{aligned}$$

Initial values

e_0	= 1,	e_{10}	= 1233816,
e_1	= 1,	e_{11}	= 6558106,
e_2	= 4,	e_{12}	= 35202448,
e_3	= 17,	e_{13}	= 190568779,
e_4	= 77,	e_{14}	= 1039296373,
e_5	= 364,	e_{15}	= 5704834700,
e_6	= 1776,	e_{16}	= 31494550253,
e_7	= 8881,	e_{17}	= 174759749005,
e_8	= 45296,	e_{18}	= 974155147162.
e_9	= 234806,		

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Equation for closed closures

$$\begin{aligned} \mathit{Clos}_0 &::= \mathcal{L}_0 \times \square \mid \mathcal{L}_1 \times \langle \mathit{Clos}_0 \rangle \mid \\ &\quad \mathcal{L}_2 \times \langle \mathit{Clos}_0, \mathit{Clos}_0 \rangle \mid \mathcal{L}_3 \times \langle \mathit{Clos}_0, \mathit{Clos}_0, \mathit{Clos}_0 \rangle \mid \dots \\ &= \bigoplus_{m \geq 0} (\mathcal{L}_m \times \mathit{Clos}_0^m) \end{aligned}$$

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Then the equation for the generating function

$$C_0(z) = \sum_{m \geq 0} L_m(z) C_0(z)^m.$$

The smallest closed closures

size	closures	total
0, 1		0
2	$\langle \lambda \underline{0}, \square \rangle$	1
3	$\langle \lambda \lambda \underline{0}, \square \rangle$ $\langle \underline{0}, \langle \lambda \underline{0}, \square \rangle \rangle$	2
4	$\langle \lambda \lambda \lambda \underline{0}, \square \rangle$ $\langle \lambda \lambda \underline{1}, \square \rangle$ $\langle \lambda (\underline{00}), \square \rangle$ $\langle \lambda \underline{0}, \langle \lambda \underline{0}, \square \rangle \rangle$ $\langle \underline{0}, \langle \lambda \lambda \underline{0}, \square \rangle \rangle$ $\langle \underline{0}, \langle \underline{0}, \langle \lambda \underline{0}, \square \rangle \rangle \rangle$	6

▶ Plain closures

Theorem

There exist

- $\bar{\rho} < \underline{\rho}$ satisfying $\bar{\rho} < \underline{\rho} < \rho_{L_\infty}$ and
- functions $\theta(n), \kappa(n)$ satisfying $\limsup_{n \rightarrow \infty} \sqrt[n]{\theta(n)} = \limsup_{n \rightarrow \infty} \sqrt[n]{\kappa(n)} = 1$

such that for sufficiently large n we have

- $\underline{\rho}^{-n} \theta(n) < c_{0,n} < \bar{\rho}^{-n} \kappa(n)$.

Sketch of the proof (case $\underline{C}_0(z)$)

- We may think of bounding $\underline{C}_0(z)$ below by

$$\underline{C}_0(z) = \sum_{m \geq 0} L_0(z) \underline{C}_0(z)^m = L_0(z) \sum_{m \geq 0} \underline{C}_0(z)^m = \frac{L_0(z)}{1 - \underline{C}_0(z)}.$$

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but this does not work.

- We approximate λ -terms by *h-shallow* λ -terms
 - ▶ terms with de Bruijn indices bounded by h .

The approximation works for $h = 153$

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- ▶ not for $h = 152$!

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- We bound $C_0(z)$ above by

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- It works !

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- Counting **plain** environments and closed closures is relatively easy
 - ▶ This is not the most realistic model.
- Counting **closed** environments and closures is challenging.
 - ▶ This is what we want.
- The implicit and infinite specification of closed closures based on closed λ -terms complicates the quantitative analysis,

Any questions !

De Bruijn indices

▸ Back

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Variables are replaced by natural numbers.

De Bruijn indices

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$(\lambda x.x)(\lambda x.\lambda y(y x))$	$(\lambda 0)(\lambda\lambda 0 1)$

Size of terms

▶ Back

- **Abstractions** λ are of size 1.
- **Applications** are of size 0.

Three methods (at least!) for counting de Bruijn indices:

- De Bruijn indices have size 0,
- De Bruijn indices have size 1
- De Bruijn indices n have size $n + 1$

	Variable size 0	Variable size 1	natural size
$\lambda(0\ 0)$	2	4	4
$\lambda\lambda 1$	2	3	4
$(\lambda 0)(\lambda\lambda(0\ 1))$	5	8	9

The h -approximations of C_0

▶ Back

$$L_0^{(h)}(z) = zL_1^{(h)}(z) + zL_0^{(h)}(z)L_0^{(h)}(z)$$

$$L_1^{(h)}(z) = zL_2^{(h)}(z) + zL_1^{(h)}(z)L_1^{(h)}(z) + z$$

$$L_2^{(h)}(z) = zL_3^{(h)}(z) + zL_2^{(h)}(z)L_2^{(h)}(z) + z + z^2$$

...

$$L_{h-1}^{(h)}(z) = zL_h^{(h)}(z) + zL_{h-1}^{(h)}(z)L_{h-1}^{(h)}(z) + z + z^2 + \dots + z^{h-1}$$

$$L_h^{(h)}(z) = zL_h^{(h)}(z) + zL_h^{(h)}(z)L_h^{(h)}(z) + z + z^2 + \dots + z^h$$