

Statistical properties of lambda terms

Maciej Bendkowski Olivier Bodini
Sergey Dovgal

CLA, Jussieu, Paris, 24/05/2018

Chapter 1. *Historical overview*

What is a closed lambda term?

Expression such as

$$\lambda x. \lambda y. (\lambda z. (\lambda x. zx)(\lambda y. zy))(xy)$$

Expression such as

$$\lambda x. \lambda y. (\lambda z. (\lambda x. zx)(\lambda y. zy))(xy)$$

- λ – abstraction

Expression such as

$$\lambda x. \lambda y. (\lambda z. (\lambda x. z @ x) (\lambda y. z @ y)) @ (x @ y)$$

- λ – abstraction
- $@$ – application

Expression such as

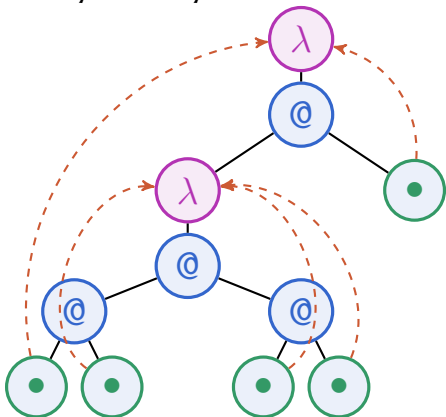
$$\lambda x. \lambda y. (\lambda z. (\lambda x. zx)(\lambda y. zy))(xy)$$

- λ – abstraction
- $@$ – application
- x, y, z – variables

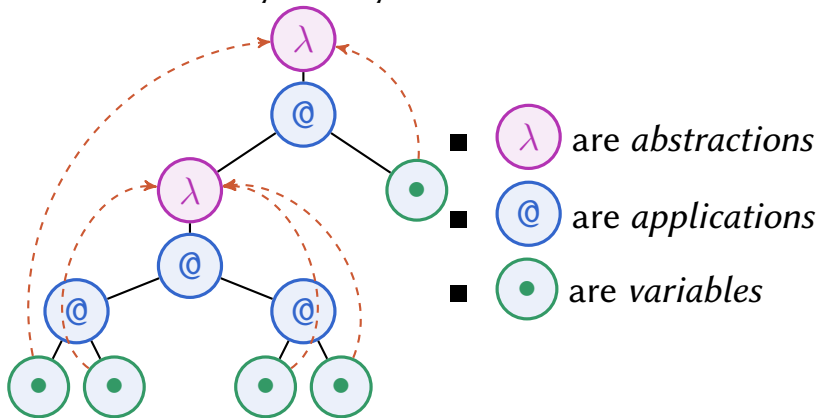
also lambda terms are..

Unary-binary trees with links

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Unary-binary trees with links



Local summary

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- TODO: What is the difference between closed and plain lambda terms?

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Motivation

First application. Software testing techniques.

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The Glasgow Haskell Compiler

		Вики	Хроно
#5557 <u>closed bug (fixed)</u>			
Code using seq has wrong strictness (too lazy)			
Сообщил:	michal.palka	Владелец:	
Приоритет:	high	Этап разработки:	
Компонент:	Compiler	Версия:	
Ключевые слова:	seq strictness strict lazy	Копия:	
Operating System:	Unknown/Multiple	Architecture:	
Type of failure:	Incorrect result at runtime	Test Case:	

Second motivation. Relation between (linear) lambda terms and (trivalent) maps.

Third motivation. Development of analytic combinatorics.

Motivation

Summary

- Software testing techniques
- Lambda terms vs. maps
- Development of analytic combinatorics

Plain and closed lambda terms

- **Closed terms:** unary-binary trees with links between variables and abstractions
- **Plain terms:** unary-binary trees with some variables unlinked

Example.

Example.

- $\lambda z.(\lambda y.zy)$ is a **closed** term

Example.

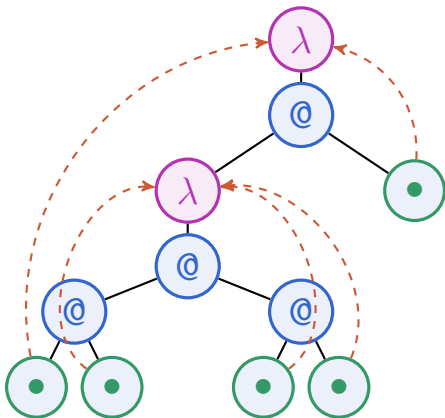
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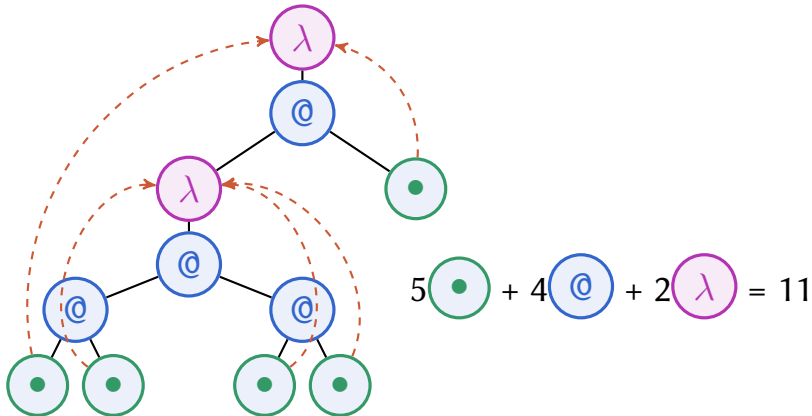
- $\lambda z.(\lambda y.zy)$ is a **closed** term
- $\lambda y.zy$ is a **plain** term
- $(\lambda x.x)(\lambda y.xy)$?

Size notion of a lambda term

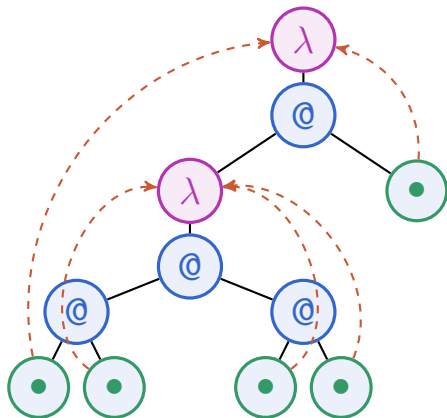
What is the size of



First approach. Total number of nodes.

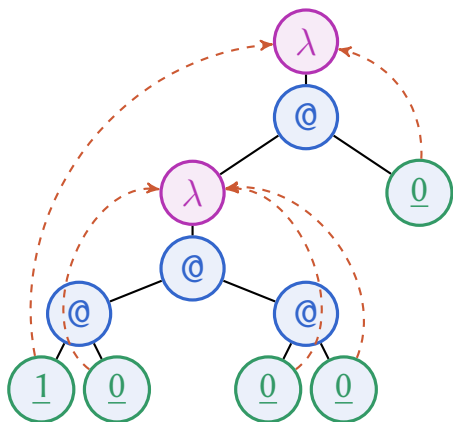


Second approach. Only abstractions and applications.



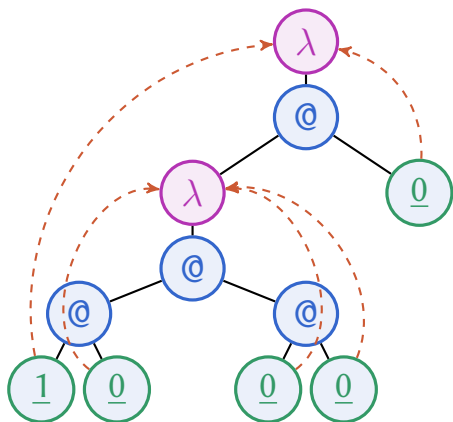
$$4 \textcircled{\text{@}} + 2 \textcircled{\lambda} = 6$$

Third approach. Natural counting.



$$4 \textcircled{\textcircled{}} + 2 \textcircled{\lambda} + (2+1+1+1+1) \textcircled{\bullet} = 12$$

Fourth approach. Generalised natural counting.



$$\begin{aligned}
 &4a \textcircled{\textcircled{}} + 2b \textcircled{\lambda} + 5c \textcircled{\bullet} + \\
 &(1+0+0+0+0)d \textcircled{\bullet} \\
 &= 4a + 2b + 5c + d
 \end{aligned}$$

Size notions, recap

- Variable size = 1
- Variable size = 0
- Variable size ≥ 1 (Natural counting)
- Generalised natural counting

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In-detail analysis of size notions

1. Variable size equal 1

Variable size equal 1

- [Bodini, Gardy, Gittenberger, Jacquot '11+]
- Linear closed lambda terms
- Upper and lower bounds for asymptotics of all closed terms
- New bijections
- Continued [Zeilberger, Giorgetti + '15+]
- :(Formal power series is divergent

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2. Variable size equal 0

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- [Grygiel, Lescanne '12]
- Enumeration and random generation
- [David, Grygiel, Kozik, Raffalli, Theyssier, Zaionc '13]
- $\Theta(\sqrt{n/\log n})$ head abstractions
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3. Natural counting

Natural counting

- Variable size = unary distance to binding λ
- [Bendkowski, Grygiel, Lescanne, Zaionc '12]
- [Bodini, Gittenberger, Gołębiewski '16]
Enumeration and asymptotics
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Why consider a simpler model?

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Main reason: advantages in random generation

Why consider a simpler model?

Marking variables imply control over expectations

Why consider a simpler model?

Another reason: parameter study

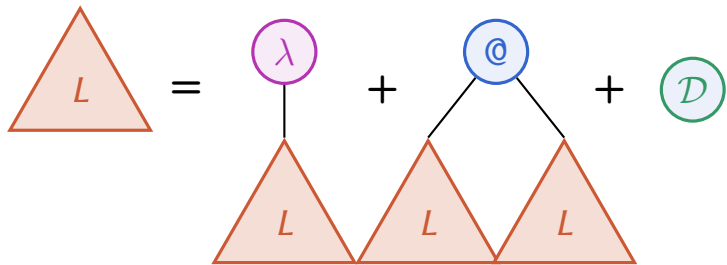
Why consider a simpler model?

Marking variables allow to study distributions

Chapter 2. *Statistical properties*

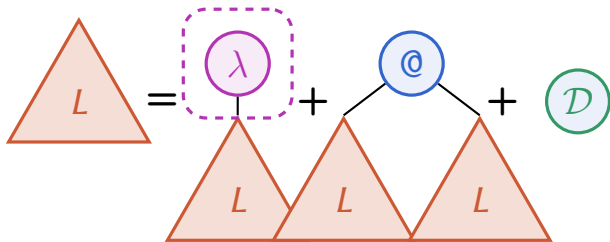
Plain lambda terms

Generating function for plain lambda terms



$$L(z) = zL(z) + zL^2(z) + \frac{z}{1-z}$$

Abstractions in plain terms?



$$L(z, u) = zuL(z, u) + zL(z, u) + \frac{z}{1-z}$$

$$L(z, u) \sim a(z, u) - b(z, u) \sqrt{1 - \frac{z}{\rho(u)}}$$

Multivariate Central Limit Theorem

Step 1. Extract coefficient of $L(z, u)$

$$L(z, u) \sim a(z, u) - b(z, u) \sqrt{1 - \frac{z}{\rho(u)}}$$

Step 2. Asymptotic behaviour of probability
generating function

$$p_n(u) \sim A(u)B(u)^n$$

Step 3. Gaussian approximation from $A(u)B(u)^n$.

Applying multivariate CLT

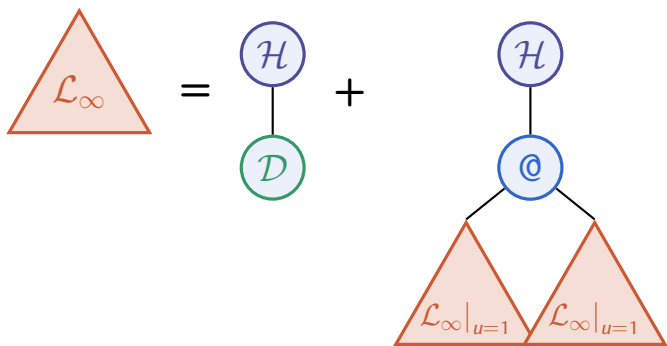
Theorem Joint Gaussian distribution for

- number of abstractions
- number of variables
- number of *redexes*

in plain lambda terms

Discrete distributions in plain terms

Head abstractions



$$L(z, u) = \frac{1}{1 - zu} \left(\frac{z}{1 - z} + zL(z, 1)^2 \right)$$

Theorem The following statistics follow discrete (geometric) limiting distributions

- The number of head abstractions
- Randomly choosed value of de Bruijn index

Chapter 3. *Infinite systems*

Drmosta–Lalley–Woods theorem

Drmotá–Lalley–Woods theorem

- Let $F(z)$ be a generating function
- Suppose it satisfies

$$F(z) = \Phi(F(z), z)$$

with Φ having combinatorial origin

- Then,

$$F(z) \sim a - b \sqrt{1 - \frac{z}{\rho}}$$

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What happens with infinite systems?

- [Drmotá, Gittenberger, Morgenbesser '12]
- *If* Jacobian is a sum of identity matrix and a compact operator
- *And* specification is strongly connected
- *Then* Infinite-dimensional version holds

Closed lambda terms satisfy an infinite system

$$L_0(z) = zL_1(z) + zL_0(z)^2 ,$$

$$L_1(z) = zL_2(z) + zL_1(z)^2 + z ,$$

$$L_2(z) = zL_3(z) + zL_2(z)^2 + z + z^2 ,$$

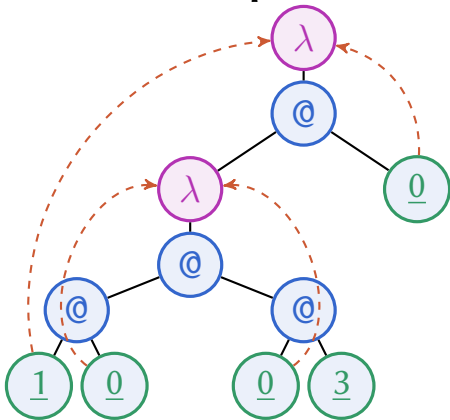
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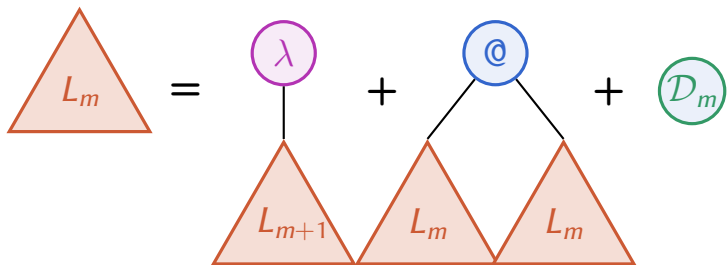
$$L_\infty(z) = zL_\infty(z) + zL_\infty(z) + \frac{z}{1-z} .$$

Why?

Adding m abstractions on the top of L_m makes the term closed

Example.

Example.



$$L_m(z) = zL_{m+1}(z) + zL_m^2(z) + z\frac{1-z^m}{1-z}$$

However

This system

$$L_0(z) = zL_1(z) + zL_0(z)^2 ,$$

$$L_1(z) = zL_2(z) + zL_1(z)^2 + z ,$$

$$L_2(z) = zL_3(z) + zL_2(z)^2 + z + z^2 ,$$

...

$$L_\infty(z) = zL_\infty(z) + zL_\infty(z) + \frac{z}{1-z} .$$

Doesn't satisfy infinite DLW theorem

Lambda terms

Infinite systems

New “master theorem”

Master theorem

Motivation

Master theorem

Motivation

$$L_m(z) = zL_{m+1}(z) + zL_m(z)^2 + z \underbrace{\frac{1 - z^m}{1 - z}}_{D_m},$$

$$L_\infty(z) = zL_\infty(z) + zL_\infty(z)^2 + z \underbrace{\frac{1}{1 - z}}_{D_\infty}.$$

The difference $D_\infty - D_m$ is exponentially small

Master theorem

More generally.

$$L_m(z) = K_m(L_m, L_{m+1}, z, u) ,$$

$$L_\infty(z) = K_\infty(L_\infty, L_\infty, z, u) ,$$

$K_\infty - K_m$ is exponentially small specified at L_∞
 L_m, K_m, u are vectors

Master theorem

Assumptions

- 1 The limiting system has Puiseux expansion
- 2 Limiting system dominates m -th system
- 3 The difference is exponentially small

Statement

- Each L_m also has Puiseux expansion*

Each L_m also has Puiseux expansion*

Two possible variants:

1 (Weak)

$$[z^n]L_m(z, u) \sim [z^n] \left(a(u) - b(u) \sqrt{1 - \frac{z}{\rho(u)}} \right)$$

2 (Strong)

The analytic continuation of $L_m(z, u)$ in a *delta-domain* satisfies

$$L_m(z, u) \sim a(u) - b(u) \sqrt{1 - \frac{z}{\rho(u)}}$$

Lambda terms

Infinite systems

Applications of “master theorem”

Lambda terms

Infinite systems

Theorem Joint Gaussian distribution for

- number of abstractions
- number of variables
- number of successors
- number of *redexes*

in closed lambda terms

Theorem Discrete distributions for

- number of head abstractions
- randomly chosen index value
- redex search time
- free variables
- missing top abstractions

in closed lambda terms

Theorem Rayleigh distributions for

- unary height profile
- natural height profile

in plain and closed lambda terms for variables,
abstractions and applications

Summary for statistics

- Normal limit laws
- New discrete limit laws
- Rayleigh limit laws

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Summary for statistics

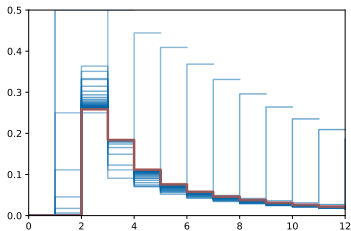
- Normal limit laws
- New discrete limit laws
- Rayleigh limit laws

Grand summary for statistics

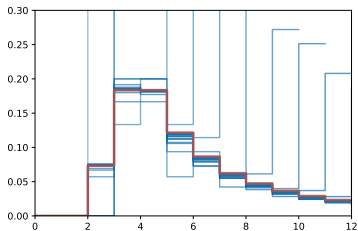
Parameter	Mean, \sim		Distribution	
	plain	closed	plain	closed
Variables	0.307 n		Normal	
Abstractions	0.258 n		Normal	
Successors	0.129 n		Normal	
Redexes	0.091 n		Normal	
Index value	0.420		Geometric	
Redex search time	6.222	6.054	Discrete	Discrete
Head abstractions	0.420	1.447	Geometric	Discrete
m -openness	2.019	0	Discrete	trivial
Free variables	5.722	0	Discrete	trivial
Unary height profile	0.122 \sqrt{n}		Rayleigh	
Natural height profile	0.412 \sqrt{n}		Rayleigh	

Epilogue. *Open problems and experiments*

Redex search time distribution

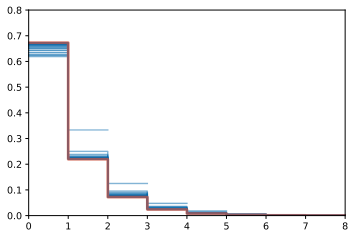


plain terms

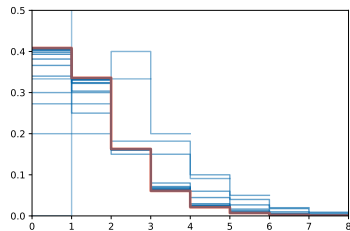


closed terms

Head abstractions

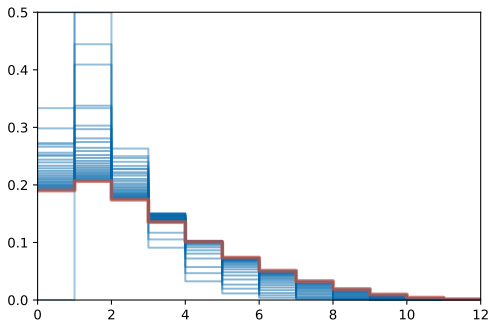


plain terms

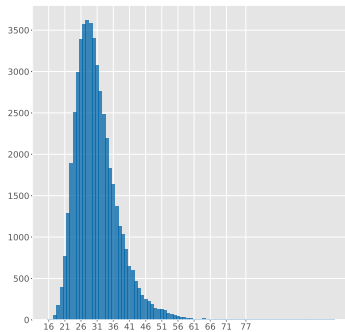


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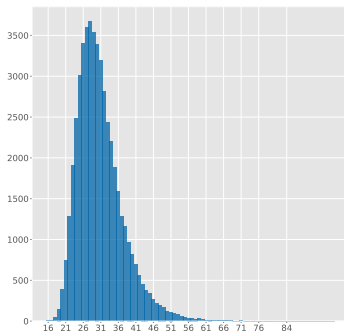
Free variables



Maximum number of variables bound to a single abstraction

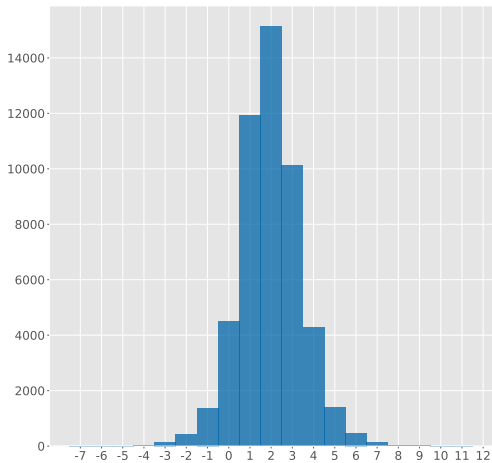


plain terms

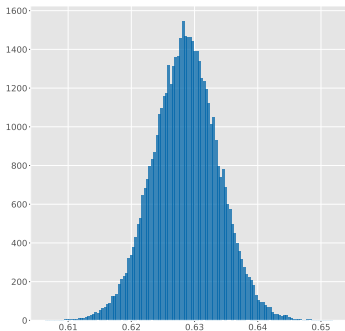


closed terms

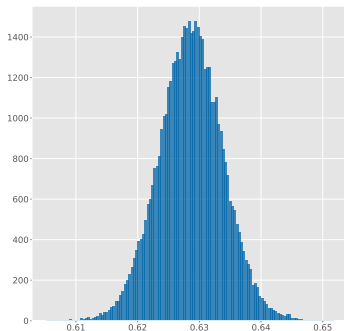
Generalised m-openness



Number of binding abstractions



plain terms



closed terms

Open questions

- Maximum number of variables bound to a single abstraction
- Number of binding abstractions
- Generalised m -openness
- Number of BCI, BCK terms in natural size notion
- Average shape of a random lambda term after k β -reductions?

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Conclusions

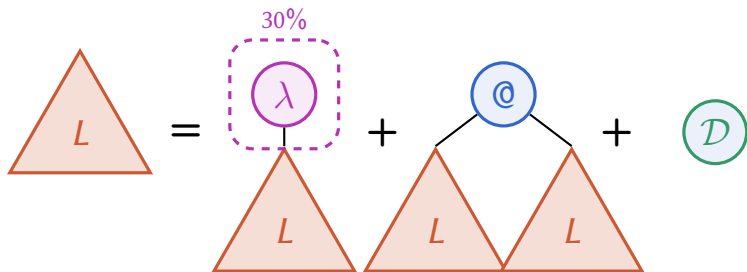
First.

First. Many tractable parameters in closed lambda terms in de Bruijn size notion

Second.

Second. Efficient multiparametric Boltzmann sampling, tweaking the desired parameters

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Third.

Third. Infinite systems of algebraic equations*

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Conjecture Premises of DLW Theorem for infinite strongly connected systems are not sufficient.

